



IBM Spectrum Scale Performance and sizing update

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ESS Packaging Options



Elastic Storage Server (ESS) is a prepacked solution using on the GNR software.¹ It comes in various models configured with different HW:

- SSD Models (400/800 GB)
 - GS1, GS2, GS4
 - 2 x High Volume Servers
 - 1/2/4 x JBOD disk enclosures
- 10,000 RPM Models (1.2 TB)
 - GS2, GS4, GS6
 - 2 x High Volume Servers
 - 2/4/6 x JBOD disk enclosures
- NL-SAS Models (2/4/6 TB)
 - GL2, GL4,GL6
 - 2 x High Volume Servers
 - 2/4/6 x JBOD disk enclosures

ESS HW Components



Servers

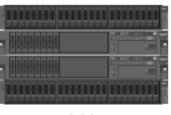


JBOD Enclosure 2U x 24, 2.5" disks



JBOD Enclosure 4U x 60, 3.5" disks

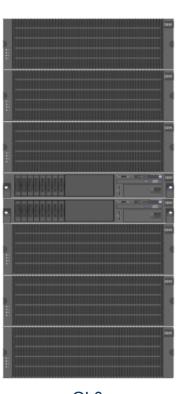
No storage controller!



GS2



GS6



GL₆

^{1.} Unlike traditional GPFS which communicates with an external block storage controller, GNR is a software storage controller that runs within GPFS, directly managing and communicating with disks.





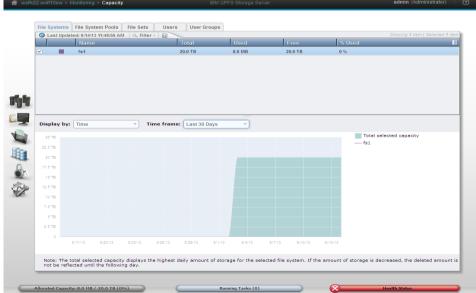
| Model | Disk size | Redundancy | Nr. Drives | Raw (TB) | Usable (TB) |
|-------|-----------|------------|------------|----------|-------------|
| GL2 | 2 | 8+2P | 116 | 232 | 170 |
| GL2 | 2 | 8+3P | 116 | 232 | 152 |
| GL2 | 4 | 8+2P | 116 | 464 | 340 |
| GL2 | 4 | 8+3P | 116 | 464 | 305 |
| GL2 | 6 | 8+2P | 116 | 696 | 510 |
| GL2 | 6 | 8+3P | 116 | 696 | 458 |
| GL4 | 2 | 8+2P | 232 | 464 | 340 |
| GL4 | 2 | 8+3P | 232 | 464 | 305 |
| GL4 | 4 | 8+2P | 232 | 928 | 680 |
| GL4 | 4 | 8+3P | 232 | 928 | 610 |
| GL4 | 6 | 8+2P | 232 | 1392 | 1020 |
| GL4 | 6 | 8+3P | 232 | 1392 | 916 |
| GL6 | 2 | 8+2P | 348 | 696 | 510 |
| GL6 | 2 | 8+3P | 348 | 696 | 458 |
| GL6 | 4 | 8+2P | 348 | 1392 | 1020 |
| GL6 | 4 | 8+3P | 348 | 1392 | 916 |
| GL6 | 6 | 8+2P | 348 | 2088 | 1530 |
| GL6 | 6 | 8+3P | 348 | 2088 | 1376 |

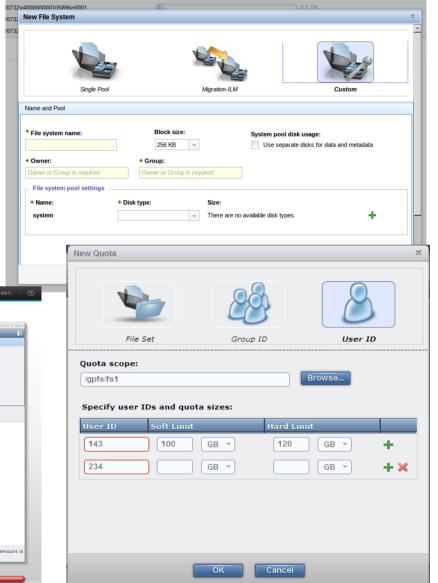
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Graphical Management



- Provide an easy-to-use Graphical User Interface for common tasks
 - System Monitoring
 - System Maintenance
 - User Configuration
- Base interface on common IBM Storage Framework
 - Comfortable for users of other IBM technologies





Quick Intro into GPFS Native Raid (GNR)



Declustered RAID

- Data and parity stripes are uniformly partitioned and distributed across a disk array.
- Arbitrary number of disks per array (unconstrained to an integral number of RAID stripe widths)
- All disks used during normal operation (no idle spares) and all disks used during rebuild

2-fault and 3-fault tolerance (RAID-D2, RAID-D3)

- Reed-Solomon parity encoding 2 or 3-fault-tolerant: stripes = 8 data strips + 2 or 3 parity strips
- 3 or 4-way mirroring

End-to-end checksum

- Disk surface to Spectrum Scale user/client
- Detects and corrects off-track and lost/dropped disk writes

Asynchronous error diagnosis while affected IOs continue on

- If media error: verify and restore if possible
- If path problem: attempt alternate paths

Advanced fault determination

- Statistical reliability and SMART monitoring
- Neighbor check, drive power cycling
- Media error detection and correction
- Slow drive detection and handling

Supports concurrent disk, enclosure and server firmware updates

Performance data



None of the following Performance numbers should be reused for sales or contract purposes.

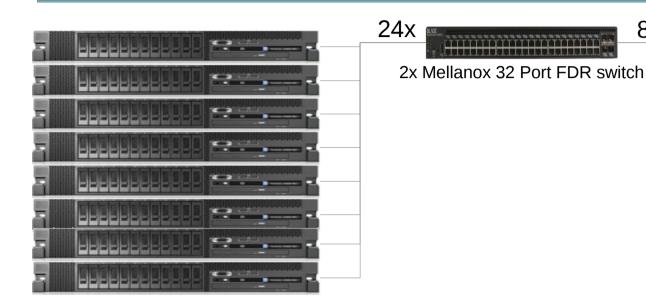
ESS Performance is typically Network bound, therefore the achievable Performance in Production depends heavily on used Network Technology and its scaling capabilities

Typical limits of Infiniband based Systems is ~25 GB/sec
Typical limits of 40GB based Systems is ~14 GB/sec
Typical limits of 10GB based Systems is ~10 GB/sec

Even if the specific ESS device is faster than above Numbers we can't guarantee the achievement of this results

Single building block Benchmark Setup





8x

12 x3650-M4 Server with 32 GB of Memory 1 FDR HBA connected 2 x 8 core CPU RHEL 7.1 GPFS 4.1.0.8 OFED-2.4-1.0.4

1 ESS GL6 System – Version 3.0 2 FDR HBA's connected per Server GPFS 4.1.0.8 code level OFED-2.4-1.0.2

Performance data IOR execution command line



Summary:

```
= POSTX
api
test filename
                   = /ibm/fs2-1m-p01/shared/ior//iorfile
                   = file-per-process
access
ordering in a file = sequential offsets
ordering inter file= no tasks offsets
clients
                   = 32 (4 per node)
repetitions
                   = 100
xfersize
                   = 1 MiB
blocksize
                 = 128 \text{ GiB}
aggregate filesize = 4096 GiB
```

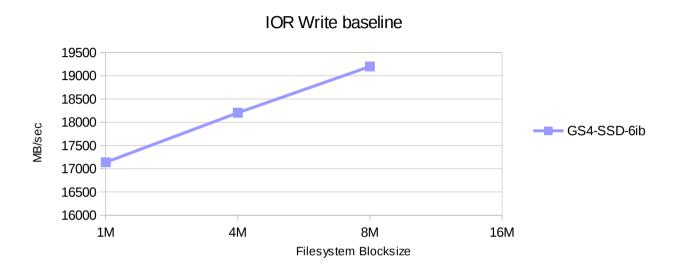


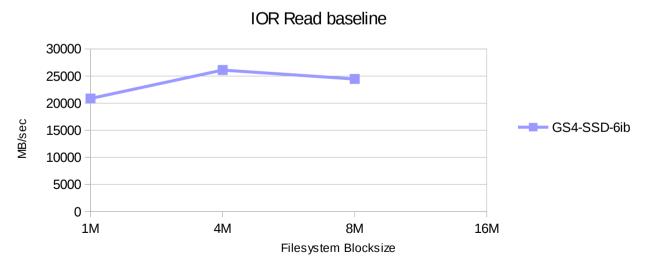


| Filesystem Blocksize | Write MB/sec | Read MB/sec |
|----------------------|--------------|-------------|
| 1 MB | 17139 | 20858 |
| 4 MB | 18205 | 26110 |
| 8 MB | 19201 | 24457 |
| 16 MB | 1 | _ |









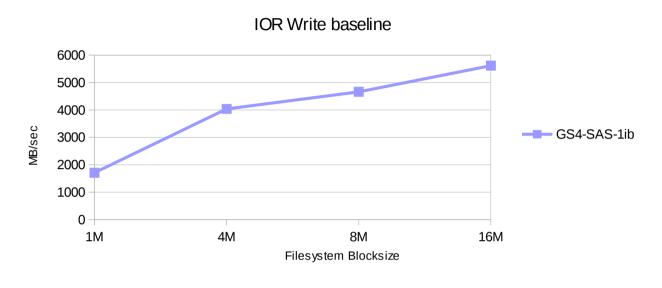


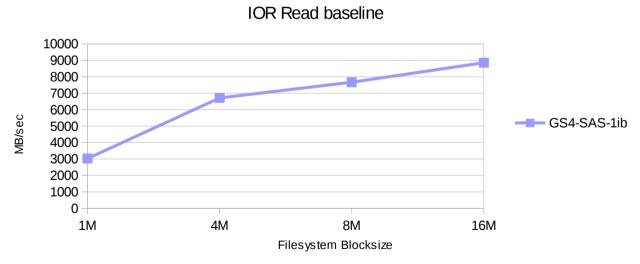


| Filesystem Blocksize | Write MB/sec | Read MB/sec |
|----------------------|--------------|-------------|
| 1 MB | 1709 | 3029 |
| 4 MB | 4039 | 6715 |
| 8 MB | 4665 | 7666 |
| 16 MB | 5619 | 8858 |









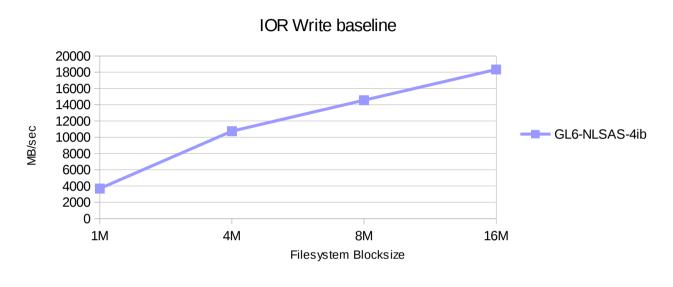


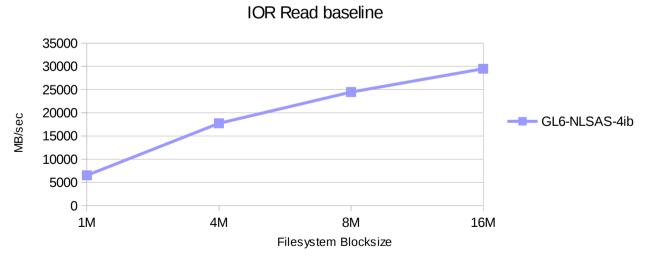


| Filesystem Blocksize | Write MB/sec | Read MB/sec |
|----------------------|--------------|-------------|
| 1 MB | 3681 | 6516 |
| 4 MB | 10748 | 17725 |
| 8 MB | 14552 | 24458 |
| 16 MB | 18337 | 29481 |











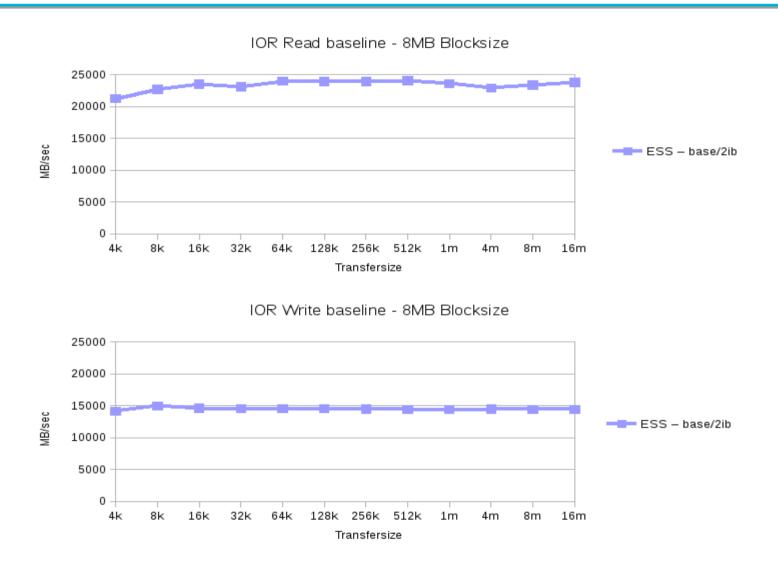


| Transfersize | Write MB/sec | Read MB/sec | | |
|--------------|--------------|-------------|--|--|
| 1 MB | 14304.18 | 23049.39 | | |
| 4 MB | 14439.93 | 23156.19 | | |
| 8 MB | 14804.21 | 23297.08 | | |
| 16 MB | 14583.56 | 21324.18 | | |

As one can see from the data, the transfersize has minimal impact on the overall throughput





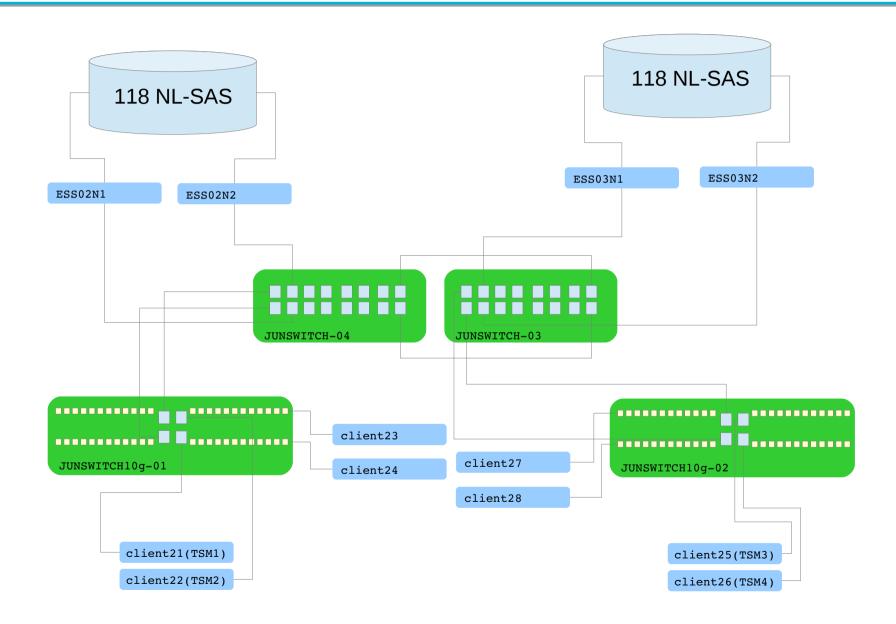




Application Specific Simulations

TSM performance testing with GL2(118 NL-SAS) *



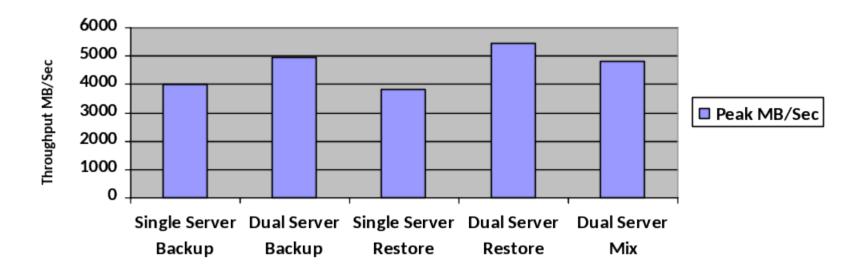


TSM performance testing with GL2(118 NL-SAS) *



Peak backup performance with multiple sessions from a single TSM server: 4017 MB/sec
Peak backup performance with multiple sessions from two TSM servers is: 4981 MB/sec
Peak restore performance with multiple sessions for a single TSM server is: 3834 MB/sec
Peak restore performance with multiple session from two TSM server is: 5424 MB/sec
Peak mixed workload performance from two TSM servers is. 4821 MB/sec

TSM and GNR Performance with 40 Gbit Ethernet



* Performance was limited by drive count in single GL2 device

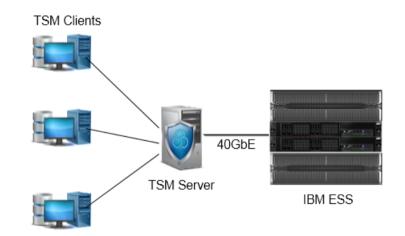
TSM Blueprint: Spectrum Protect with Elastic Storage Server GL2



Enhancements planned for 2Q15

- Support for IBM Elastic Storage Server (ESS)
 - Configuration instructions for large TSM server with ESS GL-2
 - Configuration script support for automating TSM server setup with ESS
 - Initially published for Linux x86_64
- Check https://

 ibm.biz/TivoliStorageManagerBlueprints
 for availability of the TSM Blueprint





Details of the Test Results available





Scale out backup with TSM and GSS: Performance test results

Elastic Storage with GPFS Native RAID p erformance test results with Tivoli Sto rage Manager over 40 GBit Ethernet



The peak TSM/Isilon throughput was 800MB/sec while the TSM/GPFS throughput was 5.4GB/sec (5,400MB/sec) – almost seven times faster. It's not an apples for apples comparison, but it clearly shows that Isilon is not the only fruit and GPFS could be a more flavoursome fruitstuff. With these results, acronymically TSM could stand for The Speed Machine. ®

Mirror, mirror on the wall, who has the best TSM backend of all?

Big Blue stuffs data into backup at GIGABYTES/sec

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Digital Video Simulation with Specsfs2014 VDA Benchmark



| Business | Requested | Achieved | Avg Lat | Total | Read | Write | Run | # | Cl | Avg File | Cl Data | Start Data | Init File | Max File |
|----------|-----------|----------|---------|------------|-----------|------------|-----|----|------|----------|---------|------------|-----------|-----------|
| Metric | Op Rate | Op Rate | (ms) | KBps | KBps | KBps | Sec | Cl | Proc | Size KB | Set MiB | Set MiB | Set MiB | Space MiB |
| 200 | 2000.00 | 2000.44 | 9.55 | 925240.69 | 79274.40 | 845966.29 | 300 | 10 | 40 | 1048576 | 450560 | 4505600 | 4505600 | 4915200 |
| 400 | 4000.00 | 4000.78 | 11.87 | 1843411.66 | 157015.82 | 1686395.84 | 300 | 10 | 80 | 1048576 | 901120 | 9011200 | 9011200 | 9830400 |
| 600 | 6000.00 | 6001.13 | 14.26 | 2760527.32 | 234360.60 | 2526166.72 | 300 | 10 | 120 | 1048576 | 1351680 | 13516800 | 13516800 | 14745600 |
| 800 | 8000.00 | 8001.69 | 19.44 | 3693372.00 | 313248.66 | 3380123.34 | 300 | 10 | 160 | 1048576 | 1802240 | 18022400 | 18022400 | 19660800 |
| 1000 | 10000.00 | 10001.38 | 22.63 | 4613519.11 | 394198.15 | 4219320.96 | 300 | 10 | 200 | 1048576 | 2252800 | 22528000 | 22528000 | 24576000 |
| 1200 | 12000.00 | 12001.93 | 29.69 | 5533122.30 | 471924.10 | 5061198.21 | 300 | 10 | 240 | 1048576 | 2703360 | 27033600 | 27033600 | 29491200 |
| 1400 | 14000.00 | 14001.86 | 38.63 | 6457107.05 | 550226.58 | 5906880.46 | 300 | 10 | 280 | 1048576 | 3153920 | 31539200 | 31539200 | 34406400 |
| 1600 | 16000.00 | 16000.77 | 47.30 | 7376301.04 | 630555.06 | 6745745.97 | 300 | 10 | 320 | 1048576 | 3604480 | 36044800 | 36044800 | 39321600 |
| 1800 | 18000.00 | 17999.80 | 70.28 | 8300488.87 | 709605.89 | 7590882.98 | 300 | 10 | 360 | 1048576 | 4055040 | 40550400 | 40550400 | 44236800 |
| 2000 | 20000-00 | 19517.72 | 125.49 | 9004564.53 | 789208.01 | 8215356.52 | 300 | 10 | 400 | 1048576 | 4505600 | 45056000 | 45056000 | 49152000 |

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Object Storage leveraging ESS GL6 - Setup



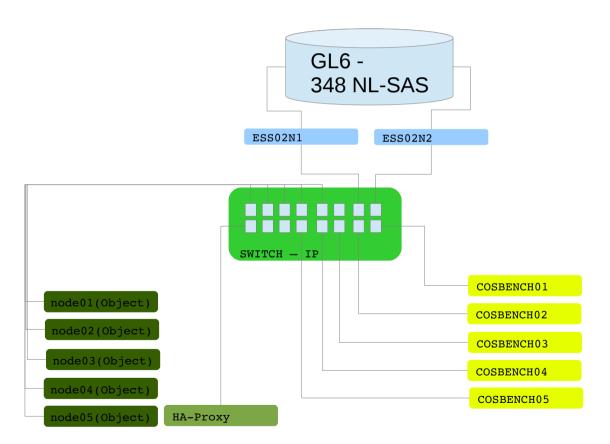
HW Setup :

Spectrum Scale Object Protocol Nodes:

5 x IBM \times 3650 M4 with 16 cores , 64 GB memory

Spectrum Scale Storage :

1 x ESS GL6



Protocol Node

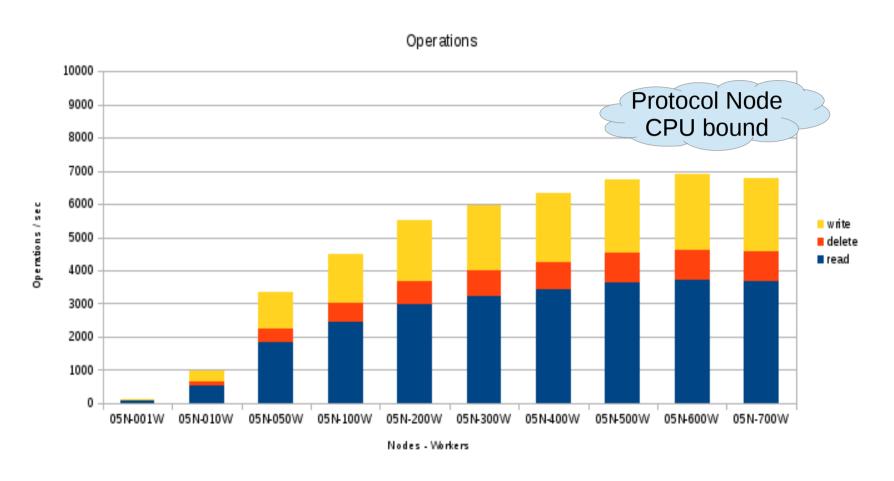
ESS Node

Workload Balancer

Workload generator







**This are preliminary numbers with the upcoming 4.1.1 release without any significant tuning effort

50 Worker Read/Write 10MB files - 1.5 GB/sec



HA Proxy Limited

ID: w320 Name: T0-W050-C10-O1K-S10M-baseline Current State: finished

Final Result

General Report

| Op-Type | Op-Count | Byte-Count | Avg-ResTime | Avg-ProcTime | Throughput | Bandwidth | Succ-Ratio |
|---------|------------|------------|-------------|--------------|-------------|-------------|------------|
| read | 29.52 kops | 295.24 GB | 337.39 ms | 47.28 ms | 98.49 op/s | 984.94 MB/S | 100% |
| delete | 7.1 kops | 0 B | 37.81 ms | 37.81 ms | 23.68 op/s | 0 B/S | 100% |
| write | 18.12 kops | 181.25 GB | 262.42 ms | 153 ms | 60.47 op/s | 604.66 MB/S | 100 |

**This are preliminary numbers with the upcoming 4.1.1 release without any significant tuning effort



Latency Tests GS4-SSD

Random 4k Read (cache Miss)



```
[root@client01 ~]# /perform/qpfsperf-mpi read rand -r 4k -n 1g -dio /ibm/fs2-1m-p01/shared/random/test-large-client01-01
/perform/gpfsperf-mpi read rand /ibm/fs2-1m-p01/shared/random/test-large-client01-01
 recSize 4K nBytes 1G fileSize 50G
 nProcesses 1 nThreadsPerProcess 1
 file cache flushed before test
 not using data shipping
 using direct I/O
 offsets accessed will cycle through the same file segment
 not using shared memory buffer
  not releasing byte-range token after open
    Data rate was 6910.39 Kbytes/sec, iops was 1727.60, thread utilization 1.000
    Record size: 4096 bytes, 1073741824 bytes to transfer, 1073741824 bytes transferred
   CPU utilization: user 1.42%, sys 1.04%, idle 97.41%, wait 0.12%
```

1727 IOPS translates to 0.579 ms / request

Seq 4k Read (cache Miss)



```
[root@client01 ~]# /perform/gpfsperf-mpi read seq -r 4k -n 1g -dio /ibm/fs2-1m-p01/shared/random/test-large-$HOSTNAME-02
/perform/qpfsperf-mpi read seq /ibm/fs2-1m-p01/shared/random/test-large-client01-02
 recSize 4K nBytes 1G fileSize 50G
  nProcesses 1 nThreadsPerProcess 1
  file cache flushed before test
  not using data shipping
  using direct I/O
  offsets accessed will cycle through the same file segment
  not using shared memory buffer
  not releasing byte-range token after open
   Data rate was 22977.24 Kbytes/sec, iops was 5744.31, thread utilization 1.000
   Record size: 4096 bytes, 1073741824 bytes to transfer, 1073741824 bytes transferred
   CPU utilization: user 1.29%, sys 1.43%, idle 97.16%, wait 0.12%
```

5744 IOPS translates to 0.174 ms / request

Seq 4k Read (cache hit)



```
[root@client01 ~]# /perform/gpfsperf-mpi read seq -r 4k -n 1g -dio /ibm/fs2-1m-p01/shared/random/test-large-$HOSTNAME-02
/perform/qpfsperf-mpi read seq /ibm/fs2-1m-p01/shared/random/test-large-client01-02
 recSize 4K nBytes 1G fileSize 50G
 nProcesses 1 nThreadsPerProcess 1
 file cache flushed before test
 not using data shipping
 using direct I/O
 offsets accessed will cycle through the same file segment
 not using shared memory buffer
 not releasing byte-range token after open
   Data rate was 39279.57 Kbytes/sec, iops was 9819.89, thread utilization 1.000
   Record size: 4096 bytes, 1073741824 bytes to transfer, 1073741824 bytes transferred
   CPU utilization: user 1.65%, sys 1.93%, idle 96.36%, wait 0.06%
```

9819 IOPS translates to 0.101 ms / request

Random 4k Write



```
[root@client01 ~]# /perform/gpfsperf-mpi write rand -r 4k -n 1g -dio /ibm/fs2-1m-p01/shared/random/test-large-$HOSTNAME-02
/perform/qpfsperf-mpi write rand /ibm/fs2-1m-p01/shared/random/test-large-client01-02
 recSize 4K nBytes 1G fileSize 50G
 nProcesses 1 nThreadsPerProcess 1
 file cache flushed before test
 not using data shipping
 using direct I/O
 offsets accessed will cycle through the same file segment
 not using shared memory buffer
 not releasing byte-range token after open
 no fsync at end of test
   Data rate was 14174.92 Kbytes/sec, iops was 3543.73, thread utilization 1.000
    Record size: 4096 bytes, 1073741824 bytes to transfer, 1073741824 bytes transferred
    CPU utilization: user 1.80%, sys 1.42%, idle 96.67%, wait 0.11%
```

3543 IOPS translates to 0.282 ms / request

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Node Sizing



| Protocol | Min Memory Recommendation | Min CPU Socket Recommendation |
|--|---|--------------------------------------|
| NFS 4000 connections per node <= 128K / cluster (32 NFS nodes * 4000) | base of >= 64GB | 1 CPU socket |
| SMB 3000 connections per node / <= 20K / cluster | x2 memory from a base of 64GB >= 128GB | 2 CPU socket |
| Object 2000 connections per node <= 32K / cluster (16 object nodes * 2000) | x2 memory from a base of 64GB >= 128GB | 2 CPU socket |
| Combination of multiple protocols 2000 Object + 3000 SMB + 4000 NFS per node | Any mix of protocols x2 memory from a base of 64GB >= 128GB | Any mix of protocols 2 CPU socket |



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Contact

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Details

Certification Date 2015-04-01 Certified Until 2018-03-31

Server System IBM Elastic Storage Server

GS2 Model

Storage Connector NAS - Distributed file system

Models included into Storage Family: IBM ESS Restrictions & GS2, GS4, GS6, GL2, GL4, GL6; Scale out up to 8 Comments

SAP HANA nodes (1TB) per GS2 (SSD); up to 16 SAP HANA nodes (1TB) per GL6 (HDD);

Connection options: 10/40 GBit Ethernet, 56 GBit Infiniband: See SAP Notes 784391 and 1084263

for GPFS support on SAP HANA nodes.

http://global.sap.com/community/ebook/2014-09-02-hana-hardware/enEN/enterprise-storage.html

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writes to sparse files, e.g., VM disk images





| Common issue in Small and medium-sized workloads EDA workload | One way traversal time at each layer | NAS Client (NFS/CIFS) Client side Network |
|---|--------------------------------------|--|
| Virtual Machine Solutions Issue across wide range of workloa | 50 us - 0.5 ms V | CTDB / NAS Layer (NFS/CIFS) |
| VMsDatabases | 50 us | Spectrum Scale FS |
| Windows home directoryLoggingISSM (ECM, Websphere, etc) | \rightarrow | NSD Client |
| o Require low-latency and non-volati | le 50 us | IB Network |
| memory Flash-backed DIMMs | ▼ | NSD Server |
| Large batteriesFast SSDs (Fusion-IO, etc) | 2 ms | Storage Controller |
| Cannot optimize data path in isola Recovery log updates occur on applica | | Disk |

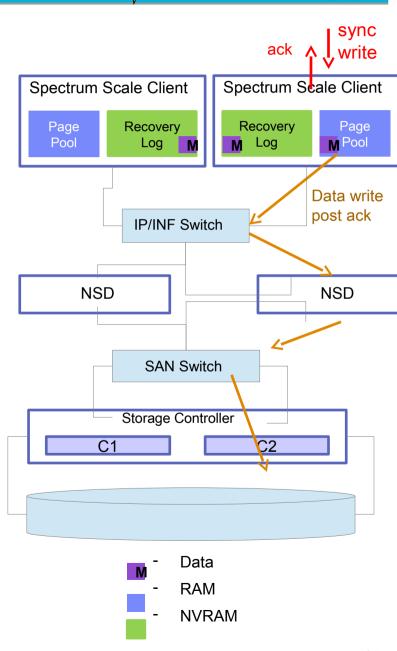
4KB Total Round Trip Time = \sim 5 ms

Solution: HAWC – Highly available Write Cache





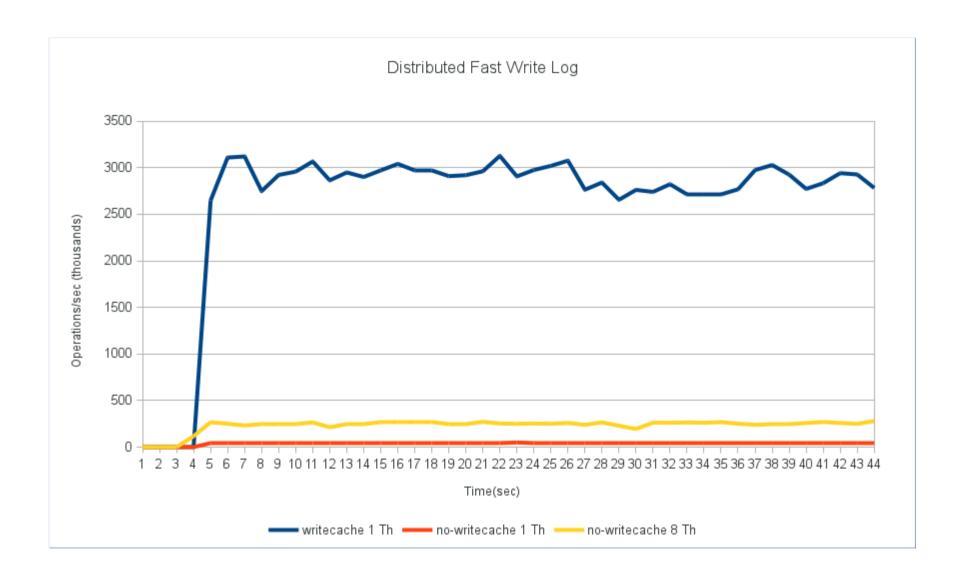
- HAWC (Log writes)
 - Store recovery log in client NVRAM
 - Either replicate in pairs or store on shared storage
 - Log writes in recovery log
 - Log small writes and send large writes directly to disk
 - Logging data only hardens it
 - Data remains in pagepool and is sent to disk post-logging
 - Leverages write gathering
 - Fast read-cache performance
 - On node failure, run recovery log to place data on disk











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